# HIGH-LEVEL PLANNING WITH EXPERIENCE GRAPHS

Aram Ebtekar, Maxim Likhachev

### Motivation

- □ **AI "dream":** oracle that solves all problems!
- Is this reasonable? Are humans oracles?
- Recent trend across Al/robotics:
  - **Learning** from **experience**, solution **reuse**
  - Next time a similar problem appears, find better solutions faster
  - Requires knowledge representation

### **Abstract Representation**

- A plan is a path from start to goal
- Graph is implicitly represented in **STRIPS** form:
  - $\square A =$ set of **atoms**: ON A B, ONTABLE B, HOLDING C
  - $\square 2^A = set of nodes or states$
  - Op = (P, A, D) is an action with sets of atoms designated as preconditions, add and delete effects, corresponding to edges: STACK A B, PICKUP C
  - **Transition**:  $(S, (P, A, D)) \rightarrow S \cup A \setminus D$  if  $S \supseteq P$
  - **D** We're at a **goal** iff  $S \supseteq G$

### Heuristic Search Planner

- Admissible relaxation: ignore delete lists
- Estimate cost to achieve **individual** atoms ■  $g_S(a) = 1 + \min_{Op} g_S(P)$  where Op adds a
  - For sets of atoms, use  $g_S(P) = \max_{p \in P} g_S(p)$
- □ Heuristic estimate to goal:  $h(S) = g_S(G)$
- Do forward weighted A\*
  - When generating S, need to compute  $g_S(a)$
  - Use dynamic programming

### **Experience Graphs**

- Originally developed for explicit graphs by SBPL
- Store edges from previously generated paths
- Inflate non E-graph edges by  $\epsilon^E$  to bias search
  - $h^{E}(S) = \min_{\pi} \sum_{i} \min\{\epsilon^{E} h(s_{i}, s_{i+1}), c^{E}(s_{i}, s_{i+1})\}$ over all **paths**  $\pi = \langle S = s_{0}, s_{1}, s_{2}, \dots, s_{N} = G \rangle$
  - $\blacksquare \epsilon h^E$  is  $\epsilon \epsilon^E$ -**consistent** provided h is consistent
- $\square$  But how can we **compute**  $h^E$  in **STRIPS**?
  - Answer: reverse Dijkstra from G on the E-graph!

## **STRIPS E-Graphs**

#### Preprocessing phase:

- Let  $N^E$  = all E-graph nodes, plus **minimal** goal state G
- **\square** Run DP to compute  $g_C(a)$  for every state  $C \in N^E$
- Now we have pairwise distance estimates  $g_C(D)$
- **\blacksquare** Reverse **Dijkstra** from G with E-graph and  $\epsilon^E g$  edges
- □ When **generating**  $S \notin N^E$ :
  - $h^{E}(S) = \min_{\pi} \sum_{i} \min\{\epsilon^{E} h(s_{i}, s_{i+1}), c^{E}(s_{i}, s_{i+1})\}$ ■ Computable by  $h^{E}(S) = \min_{C \in N^{E}} (\epsilon^{E} g_{S}(C) + h^{E}(C))$

# Analysis

## **Extensions for Future Work**

#### Generalize E-graph actions by projections

- Can "partially inflate" non E-graph edges according to some similarity measure against E-graph edges
- To a limited extent, h already acts as such a measure
- What to do when E-graph gets big?
  - "Forget" edges which have not helped recently
- Combine with other planning methods
  - Anytime incremental planning with variable-cost actions
  - Less straightforward: GRT, abstraction heuristics, etc.

### References

- Mike Phillips, Benjamin Cohen, Sachin Chitta and Maxim Likhachev, "E-Graphs: Bootstrapping Planning with Experience Graphs," Proceedings of the Robotics: Science and Systems Conference (RSS), 2012.
- Blai Bonet and Héctor Geffner, "Planning as heuristic search," Artificial Intelligence 129.1 (2001): 5-33.
- Our in-progress code based on the previous paper: <u>https://github.com/EbTech/HSP2</u>